

# Coq cheat sheet

## Notation

Propositions	Coq
$\top, \perp$	True, False
$p \wedge q$	<code>p /\ q</code>
$p \Rightarrow q$	<code>p -&gt; q</code>
$p \vee q$	<code>p \/ q</code>
$\neg p$	<code>~ p</code>
$\forall x \in A. p(x)$	<code>forall x:A, p x</code>
$\forall x, y \in A. \forall u, v \in B. q$	<code>forall (x y:A) (u v:B), q</code>
$\exists x \in A. p(x)$	<code>exists x:A, p x</code>

Sets	Coq
1	<code>unit</code>
$A \times B$	<code>prod A B</code> or <code>A * B</code>
$A + B$	<code>sum A B</code> or <code>A + B</code>
$B^A$ or $A \rightarrow B$	<code>A -&gt; B</code>
$\{x \in A \mid p(x)\}$	<code>{x:A   p x}</code>
$\sum_{x \in A} B(x)$	<code>{x:A &amp; B x}</code> or <code>sig A B</code>
$\prod_{x \in A} B(x)$	<code>forall x:A, B x</code>

Elements	Coq
$\star \in 1$	<code>tt : unit</code>
$x \mapsto f(x)$ or $\lambda x \in A. f(x)$	<code>fun (x : A) =&gt; f x</code>
$\lambda x, y \in A. \lambda u, v \in B. f(x)$	<code>fun (x y : A) (u v : B) =&gt; f x</code>
$(a, b) \in A \times B$	<code>(a,b) : A * B</code>
$\pi_1(t)$ where $t \in A \times B$	<code>fst t</code>
$\pi_2(t)$ where $t \in A \times B$	<code>snd t</code>
$\pi_1(t)$ where $t \in \sum_{x \in A} B(x)$	<code>projT1 t</code>
$\pi_2(t)$ where $t \in \sum_{x \in A} B(x)$	<code>projT2 t</code>
$\iota_1(t) \in A + B$ where $t \in A$	<code>inl t</code>
$\iota_2(t) \in A + B$ where $t \in B$	<code>inr t</code>
$t \in \{x \in A \mid p(x)\}$ because $\rho$	<code>exist t rho</code>
$\iota(t)$ where $\iota : \{x \in A \mid p(x)\} \hookrightarrow A$	<code>projT1 t</code>

## Basic tactics

When the goal is ...	... use tactic
very simple	auto, tauto or firstorder
$p \wedge q$	split
$p \vee q$	left or right
$p \rightarrow q$	intro
$\sim p$	intro
$p \leftrightarrow q$	split
an assumption	assumption
forall $x$ , $p$	intro
exists $x$ , $p$	exists $t$

To use hypothesis $H$ ...	... use tactic
$p \vee q$	destruct $H$ as [ $H_1   H_2$ ]
$p \wedge q$	destruct $H$ as [ $H_1 H_2$ ]
$p \rightarrow q$	apply $H$
$p \leftrightarrow q$	apply $H$
$\sim p$	apply $H$ or elim $H$
False	contradiction
forall $x$ , $p$	apply $H$
exists $x$ , $p$	destruct $H$ as [ $x G$ ]
$a = b$	rewrite $H$ or rewrite $\leftarrow H$

If you want to ...	... then use
prove by contradiction $p \wedge \neg p$	absurd $p$
simplify expressions	simpl
prove via intermediate goal $p$	cut $p$
prove by induction on $t$	induction $t$
pretend you are done	admit
import package $P$	Require Import $P$
compute $t$	Eval compute in $t$
print definition of $p$	Print $p$
check the type of $t$	Check $t$
search theorems about $p$	SearchAbout $p$

## Inductive definitions

### Inductive definition of $X$

```
Inductive X args :=  
  | constructor1 : args1 -> X  
  | constructor2 : args2 -> X  
  ...  
  | constructorN : argsN -> X.
```

Coq generates induction and recursion principles `X_ind`, `X_rec`, `X_rect`.

### Construction of an object by cases

```
match t with  
  | case1: result1  
  | case2: result2  
  ...  
  | caseN: resultN  
end
```

### Recursive definition of $f$

```
Fixpoint f args := ...
```